

Relational Databases

Relational Data Model:

- introduced by T. Codd (IBM) in 1970.
- simple, mathematical foundations (mathematical relation, set theory)
- implemented in a large number of commercial systems

SQL: standard for commercial RDBMSs

Previous Data Models: Hierarchical Model, Network Model

proposed in the 60s. implemented in 70s, 80s.

Relational Model Concepts

- ▶ represent a database as a collection of relations.
- ▶ each relation looks like a table of values.
- ▶ each row in the table represents a collection of related data values.
- ▶ each row in the table represents a fact that corresponds to a real-world entity or relationship
- ▶ the table name and the column name help to interpret the meaning of the values in each row.
- ▶ all values in a column are of the same data type.

Relational Model Terminology

- a row is called a **tuple**
- a column name is called an **attribute**
- a table is called a **relation**
- the data type of the values that can appear in a column is called a **domain**

tuples, attributes, relations, domains

A **Domain** D is a set of atomic values. (atomic means indivisible wrt the relational model)

To specify a domain D , we can specify a data type, from which the data values of the domain can be taken.

Example: **Some Domains**

- Canadian/US phone numbers, set of valid 11-digit phone numbers
- Canadian SIN, set of valid SIN
- Names: Smith, Lopez ...
- Employee_Age: possible ages of employees of a company
- Dept names in a University: Computer Science, Mathematics, Chemistry
...
- Dept codes in a University: CS, MATH, CHEM ...

these are logical definitions of domains.

a **data type** or **format** is also specified for each domain.

- Canadian/US phone numbers: 1(ddd)ddd-dddd
- Canadian SIN: ddd ddd ddd
- Employee_Age: integer $\in [18, 65]$

A domain is equipped with a name, data type, format.

Often, additional information is required to interpret the values of the domain (e.g. distance in km or miles)

A **relation schema** \mathbf{R} , denoted by $\mathbf{R}(A_1, A_2, \dots, A_n)$ is a relation name \mathbf{R} , and a list of attributes $\mathbf{R}(A_1, A_2, \dots, A_n)$.

Each attribute A_i has a domain $\text{dom}(A_i)$.

A relation schema describes a relation \mathbf{R} .

R is called the **name** of the relation.

The **degree** of the relation is the number of attributes **n**.

STUDENT(Name, SSN, HomePhone, Address, OfficePhone, Age, GPA)

name? attributes? degree? domains of the attributes?

null values represent attributes whose values are unknown or do not exist.

each tuple in the relation STUDENT represents a particular student.

A **relation** r of the relation schema $\mathbf{R}(A_1, A_2, \dots, A_n)$, also denoted by $r(\mathbf{R})$, is a set of n -tuples $r = \{t_1, t_2, \dots, t_m\}$.

Each n -tuple t is an ordered list of n values $t = \langle v_1, v_2, \dots, v_n \rangle$ where each value $v_i \in \text{dom}(A_i)$ for $i = 1, \dots, n$ or has the special value null.

The i^{th} value in an n -tuple t , corresp. to the attribute A_i , is referred to as $t[A_i]$.

Alternative definition of a relation

A **relation** $r(\mathbf{R})$ is a mathematical relation (subset of the cartesian product) of degree n on the domains $\text{dom}(\mathbf{A}_i)$.

$$r(\mathbf{R}) \subset \text{dom}(\mathbf{A}_1) \times \cdots \times \text{dom}(\mathbf{A}_n)$$

The cartesian product contains all possible combinations of values from the domains.

The total number of tuples in the cartesian product (provided finiteness) is

$$|\text{dom}(\mathbf{A}_1)| * \cdots * |\text{dom}(\mathbf{A}_n)|$$

The relation contains only those valid tuples that represent a particular state of the real world. As the state of the real world changes, so does the relation.

The relation schema R does not change very often. (except when we need to add a new attribute).

Several attributes can have the same domain.

Attributes specify different roles/interpretations for the same domain. (e.g. phone numbers \rightarrow HomePhone, OfficePhone)

Characteristics of relations

How does a relation differ from a file or a table?

▶ **Ordering of tuples in a relation**

Tuples in a relation don't have any particular order.

In a file or a table, there is a certain order.

Many logical orders can be specified on a relation.

▷ **Ordering of values within a tuple**

An n-tuple is an ordered list of n values.

The ordering of values in a tuple (attributes in a relation schema) is important.

The order is irrelevant, as long as the correspondence between attributes and values is known.

This points to another definition of a relation, making the order of values in a tuple unnecessary.

▷ **Values in the tuples**

Each value in a tuple is an atomic value.

Composite and multivalued attributes are not allowed.

(1st NF assumption)

Multivalued attributes: represented by extra relations

Composite attributes: represented by their components

Special role of the null value (unknown, does not apply to, unavailable)

proposition to have many different types of null values

▷ **Interpretation of a relation**

relation schema == declaration or assertion.

the **STUDENT** relation asserts that ...

each tuple in the relation represents a fact
or a particular instance of the assertion.

some relations represent facts about entities/relationships.

the relational model represents facts uniformly as relations

relation schema == predicate (**PROLOG**, deductive DBs)

values in each tuple, satisfy the predicate.

Relational Model Notation

- relation schema R of degree n , $R(A_1, A_2, \dots, A_n)$
- an n -tuple of a relation $r(R)$, $t = \langle v_1, v_2, \dots, v_n \rangle$
- $t[A_i]$ and $t.A_i$ refer to the value v_i in t of attribute A_i
- $t[A_{i_1}, A_{i_2}, \dots, A_{i_k}]$ and $t.(A_{i_1}, A_{i_2}, \dots, A_{i_k})$
(where $A_{i_1}, A_{i_2}, \dots, A_{i_k}$ is a list of k attributes from R)
refer to the subtuple of values $t = \langle v_{i_1}, v_{i_2}, \dots, v_{i_k} \rangle$ in t corresponding to the k attributes
- Q, R, S denote relation names
- q, r, s denote relation states
- t, u, v denote tuples
- the name of a relation schema, also indicates the current set of tuples in the relation (relation state)

- an attribute A can be prefixed by the relation name R in which it belongs using the dot notation $R.A$

e.g. $STUDENT.Name$, $STUDENT.SIN$

- The same name can be used for 2 attributes in different relations
- All attribute names within one relation must be unique.